# Fast and Precise: Adjusting Planning Horizon with Adaptive Subgoal Search

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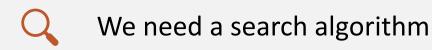
PRESENTED BY MICHAŁ TYROLSKI



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Let say that we hale "hard" environment

Given a problem P, we would like to solve it "efficiently"





It should be efficient in terms of computional budget

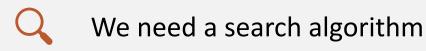
### Motivation



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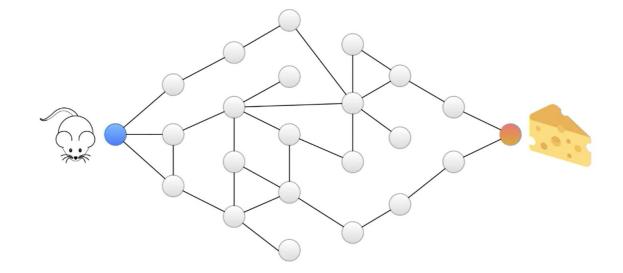




It should be **efficient** in terms of **computional budget** 

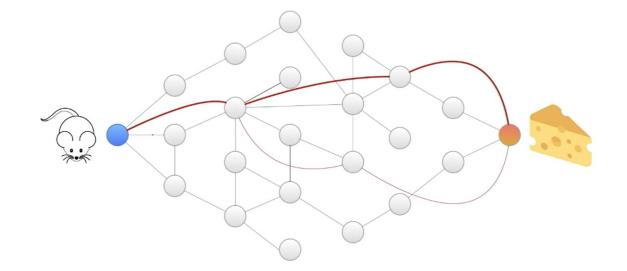
### Motivation

### Subgoal Search



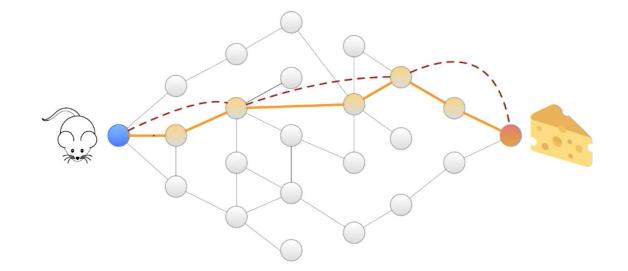
Czechowski, K., Odrzygóźdź, T., Zbysiński, M., Zawalski, M., Olejnik, K., Wu, Y., Kuciński, Ł. and Miłoś, P., 2021. Subgoal search for complex reasoning tasks. *Advances in Neural Information Processing Systems*, *34*, pp.624-638.

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Cost is defined as a number of nodes which at least one were inferenced by a neural network.



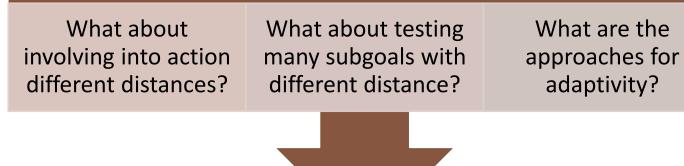
Computional budget is defined as a maximum allowed cost. If search doesn't finish after it with found solution, we say that problem is not solved.

Efficient approach menas that our search is able to achieve "good" solve rate given small computional budget

### Definitions

# This solution is great but...

#### K (distance between subgoals) is fixed



Filling gaps between subgoals with policy is very **costly**, especially for bigger K

Do we hale always to fill those gaps?

### Adaptivity: Motivational example



### Adaptive approaches

**MixSubS**: for which node in search graph, we generate on subgoal of each possible distance.

In each iteration, MixSubS chooses a state with the highest value estimation V (s) to process

**Strongest-first** uses one generator at a time with the longest distance not previously used in state s. In each iteration, Strongest-first chooses a state with the highest value estimation V (s) to process.

**Iterative Mixing** is similar to MixSubS and allows for advanced schedules of generators to be used. In the consecutive iterations, the i-th generator is used to expand l\_i nodes before switching to the next generator. This allows us to flexibly prioritize the better generators, but at the cost of tuning additional hyperparameters l1, . . . , In. For these reasons, it is not practical, but useful as a reference point.

**Longest-first** prioritizes long subgoals over the whole search procedure. Only if the queue does not contain any nodes with the higher k, it uses subgoals of lower distances. The nodes are processed in the order of their value estimation V (s).



## Filling gaps

Filling gaps between subgoals is costly.

The longest distance to subgoal, the quality of Subgoal Generator gets much worse

Weak Subgoal Generator is able to generate many invalid subgoals

Running policy for filling those gaps has no sense – we would like to avoid it

### Verifier





Verifier:: State x State  $\rightarrow$  [0, 1]



Verifier(start\_state, proposed\_subgoal\_state) predicts if policy is able to reach subgoal (proposed by subgoal generator) strating from start state and leading to subgoal state

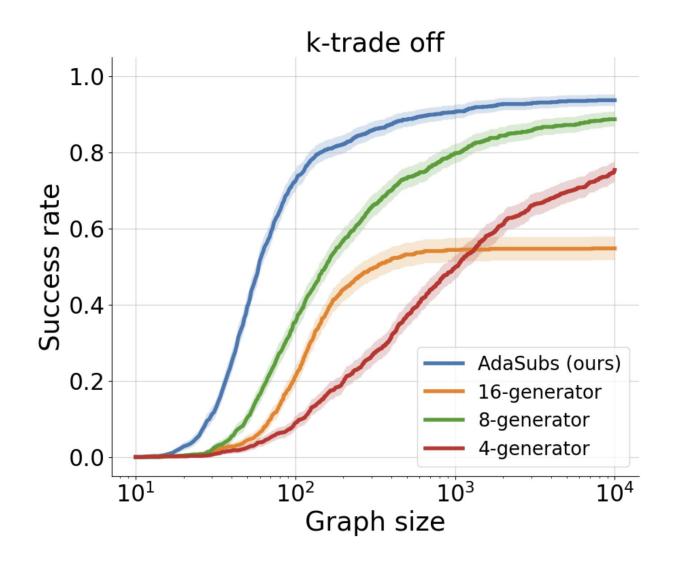


It saves a lot of computations, especially on small budgets

Algorithm 2	Condi	tional low-level policy					
<b>Requires:</b>	$C_2$	<b>^</b>					
	$\pi$						
		policy network					
	M	model of the environment					
function G	ET_PA	$ATH(s_0, subgoal)$					
$\texttt{step} \leftarrow$	- 0						
$\mathtt{s} \gets \mathtt{s}_0$							
action	_pat1	$h \leftarrow []$					
while s	tep <	$< C_2$ do					
action $\leftarrow \pi$ .PREDICT(s, subgoal)							
action_path.APPEND(action)							
$\mathtt{s} \leftarrow$	- <i>M</i> .N	EXT_STATE(s, action)					
if s	= su	bgoal <b>then</b>					
	returi	n action_path					
ste	$\mathfrak{p} \leftarrow \mathfrak{s}$	step $+1$					
return	-	<b>1</b>					

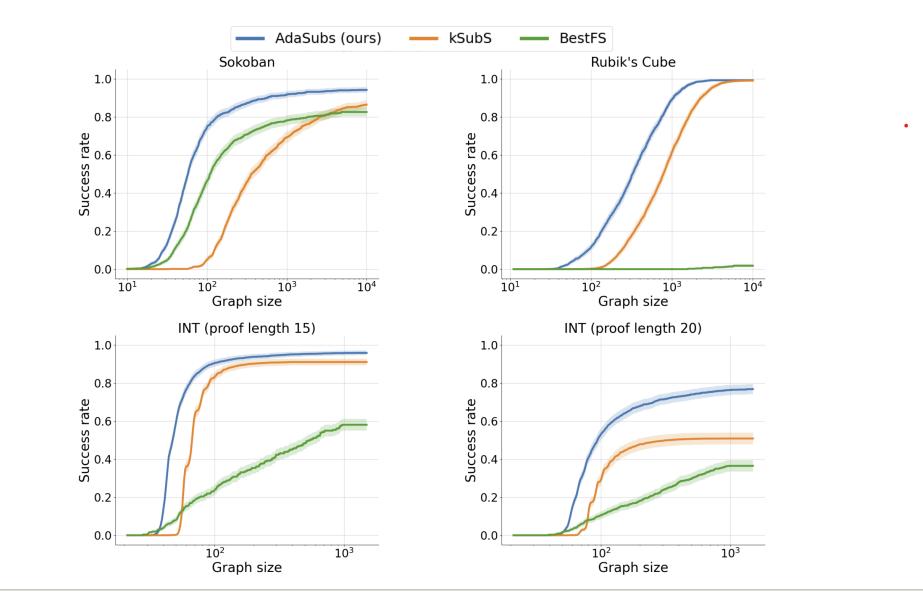
Algorithm 3 Verification algorithm							
<b>Requires:</b>	v	verifier network					
	$\mathtt{t}_{\mathtt{hi}}$	upper threshold					
	$t_{lo}$	lower threshold					
function IS_VALID( $s, s'$ )							
if $v(s, s') > t_{hi}$ then return True							
else if $v(s, s') < t_{lo}$ then return False							
return GET_PATH $(\mathbf{s}, \mathbf{s}') \neq []$							

### Results



Comparison of success rates for different subgoal generators for Sokoban. AdaSubS-k describes using a single generator with distance k.

		INT				
		Small budget (50 nodes)		Large budget (1000 nodes)		
		with verifier	without	with verifier	without	
BestFS		-	1.7%	-	36.7%	
	k = 4 $k = 3$	$2.2\%\ 4.0\%$	$0.1\%\ 0.2\%$	$82.4\%\ 89.6\%$	$83.0\%\ 90.7\%$	
kSubS	k = 2 $k = 1$	$2.1\%\ 0.0\%$	$0.5\%\ 0.0\%$	$89.8\%\ 34.7\%$	$91.7\%\ 46.0\%$	
MixSubS	$egin{aligned} k &= [4,3,2] \ k &= [3,2,1] \ k &= [3,2] \end{aligned}$	$\begin{array}{c} 0.0\% \\ 0.0\% \\ 17.0\% \end{array}$	$0.0\%\ 0.0\%\ 14.8\%$	$94.6\%\ 92.2\%\ 92.2\%$	$95.0\%\ 92.9\%\ 93.5\%$	Comparison
Iterative mixin	iterations = $[1, 1, 1]$ ng iterations = $[10, 1, 1]$ iterations = $[4, 2, 1]$	$32.0\%\ 43.0\%\ 54.0\%$	$30.1\%\ 44.8\%\ 52.1\%$	$87.0\%\ 95.1\%\ 93.6\%$	$88.6\%\ 96.0\%\ 95.5\%$	
Strongest-first	t	39.5%	40.8%	88.5%	89.8%	
Longest-first		59.0%	51.5%	95.7%	95.5%	







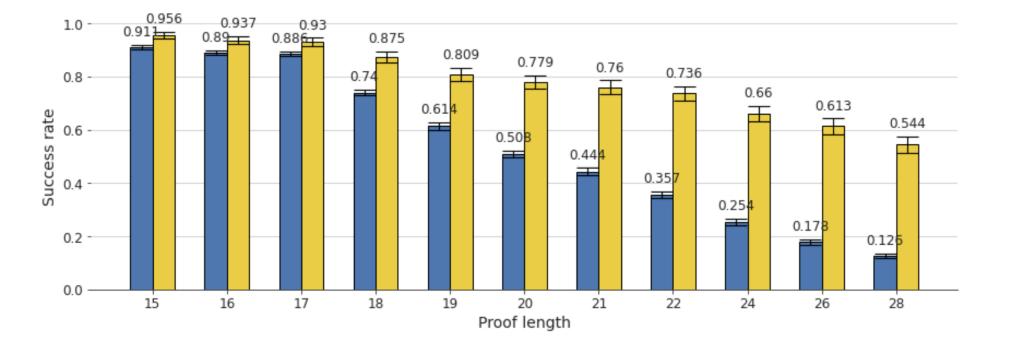


Figure 3: Out-of-distribution performance of AdaSubS and kSubS for long proofs in INT with budget of 5000 nodes. Both methods were trained on proofs of length 15. Error bars correspond to 95% confidence intervals.

Thank you!

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